

## 1 Outline

In this lecture, we study

- split cuts,
- Gomory's mixed-integer cuts,
- relationships between Gomory's mixed-integer cuts, Gomory's fractional cuts, and split cuts.

## 2 Split cuts

Consider a mixed-integer program given as follows.

$$\begin{aligned} \min \quad & c^\top x + h^\top y \\ \text{s.t.} \quad & Ax + Gy \leq b \\ & x \in \mathbb{Z}^d, y \in \mathbb{R}^p \end{aligned}$$

where  $A, G, b$  have rational entries. Then

$$P = \{(x, y) \in \mathbb{R}^d \times \mathbb{R}^p : Ax + Gy \leq b\}$$

is a rational polyhedron and the feasible set of the LP relaxation. Moreover,

$$S = P \cap (\mathbb{Z}^d \times \mathbb{R}^p)$$

is the set of solutions to the mixed-integer program.

Let  $\pi \in \mathbb{Z}^d$  and  $\pi_0 \in \mathbb{Z}$ . Note that for any  $x \in \mathbb{Z}^d$ , we know that  $\pi^\top x$  is an integer. Therefore, any  $x \in \mathbb{Z}^d$  satisfies  $\pi^\top x \leq \pi_0$  or  $\pi^\top x \geq \pi_0 + 1$ . Note that  $\mathbb{Z}^d \times \mathbb{R}^p$  can be partitioned as

$$\mathbb{Z}^d \times \mathbb{R}^p = \{(x, y) \in \mathbb{Z}^d \times \mathbb{R}^p : \pi^\top x \leq \pi_0\} \cup \{(x, y) \in \mathbb{Z}^d \times \mathbb{R}^p : \pi^\top x \geq \pi_0 + 1\}.$$

Similarly, it follows that

$$S = \{(x, y) \in S : \pi^\top x \leq \pi_0\} \cup \{(x, y) \in S : \pi^\top x \geq \pi_0 + 1\}.$$

Motivated by this, we define two polyhedron as follows.

$$\begin{aligned} \Pi_1 &= \{(x, y) \in P : \pi^\top x \leq \pi_0\}, \\ \Pi_2 &= \{(x, y) \in P : \pi^\top x \geq \pi_0 + 1\}. \end{aligned}$$

Note that

$$S \subseteq \Pi_1 \cup \Pi_2.$$

Here,  $\Pi_1 \cup \Pi_2$  is a subset of  $P$  as shown in Figure 17.1, and therefore,  $\Pi_1 \cup \Pi_2$  is a stronger relaxation of  $S$  than  $P$ .

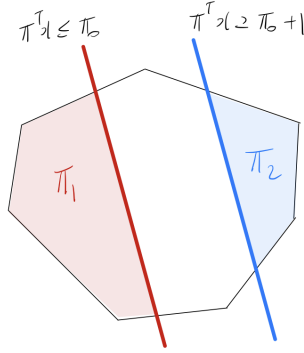


Figure 17.1: Splitting a polyhedron

We refer to an inequality  $\alpha^\top x + \beta^\top y \leq \gamma$  that is valid for  $\Pi_1 \cup \Pi_2$  as a **split cut** for  $P$ . As shown in Figure 17.2, a split cut may cut off some part of the polyhedron  $P$ .

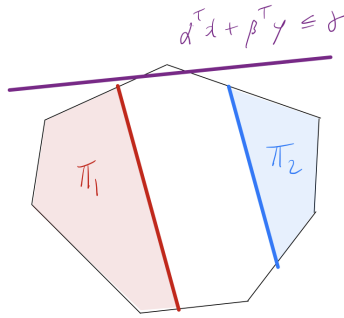


Figure 17.2: Split cut for polyhedron  $P$

The set between two parallel hyperplanes, given by

$$\{(x, y) \in \mathbb{R}^d \times \mathbb{R}^p : \pi_0 \leq \pi^\top x \leq \pi_0 + 1\}$$

is called a **split set**. We refer to  $\Pi_1 \cup \Pi_2$  obtained from the polyhedron  $P$  and  $(\pi, \pi_0) \in \mathbb{Z}^d \times \mathbb{Z}$  as a **split**. Note that splits are defined with the integer variables, not with continuous variables.

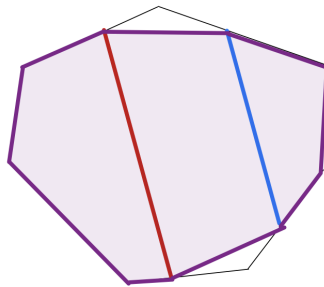


Figure 17.3: Applying all splits cuts for a split

If we apply all split cuts from the split with  $(\pi, \pi_0)$ , then we obtain  $\text{conv}(\Pi_1 \cup \Pi_2)$  as in Figure 17.3.

Recall that an inequality  $\pi^\top x \leq \pi_0$  is a Chvátal-Gomory cut if

$$P \cap \left\{ (x, y) \in \mathbb{R}^d \times \mathbb{R}^p : \pi^\top x \geq \pi_0 + 1 \right\} = \emptyset.$$

In fact, a Chvátal-Gomory cut is a split cut when one of  $\Pi_1$  and  $\Pi_2$  is empty (see Figure 17.4).

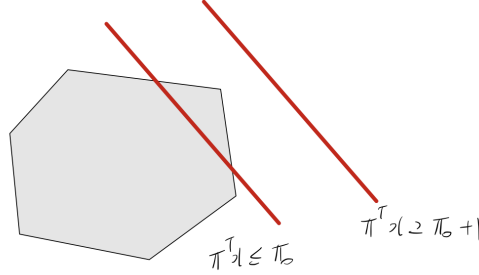


Figure 17.4: Chvátal-Gomory cut as a split cut

We may take all possible split cuts from every possible  $(\pi, \pi_0) \in \mathbb{Z}^d \times \mathbb{Z}$ . The resulting set is called the **split closure** of the polyhedron  $P$ . The split closure of  $P$  is given by

$$P^{(1)} = \bigcap_{(\pi, \pi_0) \in \mathbb{Z}^d \times \mathbb{Z}} \text{conv} \left( \Pi_1^{(\pi, \pi_0)} \cup \Pi_2^{(\pi, \pi_0)} \right)$$

where  $\Pi_1^{(\pi, \pi_0)} \cup \Pi_2^{(\pi, \pi_0)}$  denotes the split associated with  $(\pi, \pi_0) \in \mathbb{Z}^d \times \mathbb{Z}$ .

**Theorem 17.1** (Cook, Kannan, and Schrijver [CKS90]). *The split closure of any rational polyhedron is a rational polyhedron.*

Then the split closure  $P^{(1)}$  is a rational polyhedron. Therefore, we may recursively apply the procedure of taking the split closure. Let  $P^{(k)}$  denote the  $k$ th **split closure** of  $P$ , that is, the split closure of  $P^{(k-1)}$ .

**Example 17.2.** Consider a polyhedron

$$P = \left\{ (x_1, x_2, y) \in \mathbb{R}_+^3 : x_1 \geq y, x_2 \geq y, x_1 + x_2 + 2y \leq 2 \right\}$$

and the associated mixed-integer set

$$S = P \cap (\mathbb{Z}^2 \times \mathbb{R}).$$

Note that  $P$  is a convex combination of 4 points as follows.

$$P = \text{conv} \left( \left\{ (0, 0, 0), (2, 0, 0), (0, 2, 0), \left( \frac{1}{2}, \frac{1}{2}, \frac{1}{2} \right) \right\} \right).$$

Here,  $P$  has an apex  $(1/2, 1/2, 1/2)$ . Moreover,

$$\begin{aligned} S &= \{(0, 0, 0), (2, 0, 0), (0, 2, 0)\} \\ &= \{(x_1, x_2, y) \in \mathbb{R}_+^3 : x_1 \geq y, x_2 \geq y, y \leq 0\}. \end{aligned}$$

Therefore, to obtain the convex hull of  $S$ , we need to deduce inequality  $y \leq 0$ . However, we can argue that the  $k$ th split closure of  $P^{(k)}$  for any finite  $k$  contains a point of the form  $(1/2, 1/2, t)$  for some  $t > 0$ .

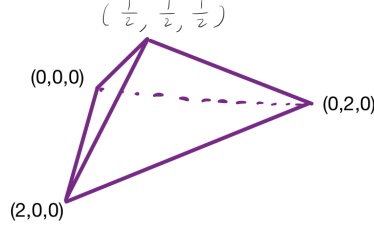


Figure 17.5:  $P$  as the convex hull of some points

### 3 Gomory's mixed-integer cuts

We consider the following mixed-integer set

$$S = \left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]} a_i x_i + \sum_{j \in [p]} g_j y_j = b \right\} \subseteq \mathbb{Z}_+^d \times \mathbb{R}_+^p.$$

Let  $f_0, f_1, \dots, f_d$  be defined as

$$f_0 = b - \lfloor b \rfloor, \quad f_i = a_i - \lfloor a_i \rfloor \quad \text{for } i \in [d].$$

We assume that

$$0 < f_0 < 1,$$

i.e.,  $b$  is not an integer, so that we may generate a nontrivial cut. Then

$$\sum_{i \in [d]} a_i x_i + \sum_{j \in [p]} g_j y_j = b$$

is equivalent to

$$\sum_{i \in [d]} f_i x_i + \sum_{j \in [p]} g_j y_j = f_0 + \left( \lfloor b \rfloor - \sum_{i \in [d]} \lfloor a_i \rfloor x_i \right).$$

This is further reduced to

$$\sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j = f_0 + \left( \lfloor b \rfloor - \sum_{i \in [d]: f_i \leq f_0} \lfloor a_i \rfloor x_i - \sum_{i \in [d]: f_i > f_0} (\lfloor a_i \rfloor + 1) x_i \right).$$

Therefore, we deduce that

$$S \subseteq \left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j = f_0 + k \quad \text{for some integer } k \right\}.$$

The set on the right-hand side is contained in

$$\left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j \geq f_0 \right\} \\ \cup \left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j \leq f_0 - 1 \right\},$$

which is equivalent to

$$\left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]: f_i \leq f_0} \frac{f_i}{f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{f_i - 1}{f_0} x_i + \sum_{j \in [p]} \frac{g_j}{f_0} y_j \geq 1 \right\} \\ \cup \left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]: f_i \leq f_0} \frac{-f_i}{1 - f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{1 - f_i}{1 - f_0} x_i + \sum_{j \in [p]} \frac{-g_j}{1 - f_0} y_j \geq 1 \right\},$$

Then

$$\sum_{i \in [d]: f_i \leq f_0} \max \left\{ \frac{f_i}{f_0}, \frac{-f_i}{1 - f_0} \right\} x_i + \sum_{i \in [d]: f_i > f_0} \max \left\{ \frac{f_i - 1}{f_0}, \frac{1 - f_i}{1 - f_0} \right\} x_i + \sum_{j \in [p]} \max \left\{ \frac{g_j}{f_0}, \frac{-g_j}{1 - f_0} \right\} y_j \geq 1$$

is a valid inequality for  $S$ . Since  $0 \leq f_i < 1$  for  $i \in [p]$ . Here, the inequality is equal to

$$\sum_{i \in [d]: f_i \leq f_0} \frac{f_i}{f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{1 - f_i}{1 - f_0} x_i + \sum_{j \in [p]: g_j \geq 0} \frac{g_j}{f_0} y_j + \sum_{j \in [p]: g_j < 0} \frac{-g_j}{1 - f_0} y_j \geq 1.$$

This inequality is a **Gomory's mixed-integer (GMI) cut**.

### 3.1 Comparison with Gomory's fractional cuts

Recall that we considered sets of the form

$$S = \left\{ x \in \mathbb{Z}_+^d : \sum_{i \in [d]} a_i x_i = b \right\}.$$

Recall that Gomory's fractional cut is given by

$$\sum_{i \in [d]} (a_i - [a_i]) x_i \geq b - [b].$$

Then Gomory's fractional cut is equivalent to

$$\sum_{i \in [d]} \frac{f_i}{f_0} x_i \geq 1$$

where  $f_i = a_i - [a_i]$  for  $i \in [d]$  and  $f_0 = b - [b]$ . On the other hand, Gomory's mixed-integer cut for the set  $S$  has the form

$$\sum_{i \in [d]: f_i \leq f_0} \frac{f_i}{f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{1 - f_i}{1 - f_0} x_i \geq 1.$$

In fact, if  $f_i > f_0$ , then

$$\frac{f_i}{f_0} > \frac{1 - f_i}{1 - f_0}.$$

This indicates that Gomory's mixed-integer cuts dominate Gomory's fractional cuts.

### 3.2 Connection to split cuts

Let us consider again the mixed-integer set given by

$$S = \left\{ (x, y) \in \mathbb{Z}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]} a_i x_i + \sum_{j \in [p]} g_j y_j = b \right\} \subseteq \mathbb{Z}_+^d \times \mathbb{R}_+^p.$$

Let  $(\pi, \pi_0) \in \mathbb{Z}^d \times \mathbb{Z}$  be defined as follows.

$$\pi_i = \begin{cases} \lfloor a_i \rfloor, & \text{if } f_i \leq f_0 \\ \lfloor a_i \rfloor + 1, & \text{if } f_i > f_0 \end{cases}, \quad \pi_0 = \lfloor b \rfloor.$$

Then  $(\pi, \pi_0)$  defines split  $\Pi_1 \cup \Pi_2$  given by

$$\Pi_1 = \left\{ (x, y) \in \mathbb{R}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]} a_i x_i + \sum_{j \in [p]} g_j y_j = b, \quad \sum_{i \in [d]: f_i \leq f_0} \lfloor a_i \rfloor x_i + \sum_{i \in [d]: f_i > f_0} (\lfloor a_i \rfloor + 1) x_i \leq \lfloor b \rfloor \right\},$$

$$\Pi_2 = \left\{ (x, y) \in \mathbb{R}_+^d \times \mathbb{R}_+^p : \sum_{i \in [d]} a_i x_i + \sum_{j \in [p]} g_j y_j = b, \quad \sum_{i \in [d]: f_i \leq f_0} \lfloor a_i \rfloor x_i + \sum_{i \in [d]: f_i > f_0} (\lfloor a_i \rfloor + 1) x_i \geq \lfloor b \rfloor + 1 \right\}$$

For  $\Pi_1$ , subtracting the inequality from the equality, we deduce that

$$\sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j \geq f_0$$

is valid for  $\Pi_1$ . For  $\Pi_2$ , subtracting the inequality from the equality, we deduce that

$$\sum_{i \in [d]: f_i \leq f_0} f_i x_i + \sum_{i \in [d]: f_i > f_0} (f_i - 1) x_i + \sum_{j \in [p]} g_j y_j \leq f_0 - 1$$

is valid for  $\Pi_2$ . The inequalities are equivalent to

$$\sum_{i \in [d]: f_i \leq f_0} \frac{f_i}{f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{f_i - 1}{f_0} x_i + \sum_{j \in [p]} \frac{g_j}{f_0} y_j \geq 1$$

$$\sum_{i \in [d]: f_i \leq f_0} \frac{-f_i}{1 - f_0} x_i + \sum_{i \in [d]: f_i > f_0} \frac{1 - f_i}{1 - f_0} x_i + \sum_{j \in [p]} \frac{-g_j}{1 - f_0} y_j \geq 1$$

where the first inequality is valid for  $\Pi_1$  and the second inequality is valid for  $\Pi_2$ . Then we can argue that

$$\sum_{i \in [d]: f_i \leq f_0} \max \left\{ \frac{f_i}{f_0}, \frac{-f_i}{1 - f_0} \right\} x_i + \sum_{i \in [d]: f_i > f_0} \max \left\{ \frac{f_i - 1}{f_0}, \frac{1 - f_i}{1 - f_0} \right\} x_i + \sum_{j \in [p]} \max \left\{ \frac{g_j}{f_0}, \frac{-g_j}{1 - f_0} \right\} y_j \geq 1$$

is valid for  $\Pi_1 \cup \Pi_2$ . In fact, the inequality is precisely Gomory's mixed-integer cut for the set  $S$ .

### 3.3 Gomory's mixed-integer closure

Let us consider a rational polyhedron given by

$$P = \left\{ (x, y) \in \mathbb{R}_+^d \times \mathbb{R}_+^p : Ax + Gy \leq b \right\}.$$

By adding nonnegative slack variables  $s$ , we deduce that

$$P = \left\{ (x, y) \in \mathbb{R}_+^d \times \mathbb{R}_+^p : \exists s \in \mathbb{R}_+^m \text{ s.t. } Ax + Gy + s = b \right\}$$

where  $m$  is the number of inequalities in the system  $Ax + Gy \leq b$ . Then we may take any linear combination of the equality constraints in  $Ax + Gy + s = b$  by taking a multiplier  $\lambda$ :

$$(A^\top \lambda)x + (G^\top \lambda)y + \lambda^\top s = b^\top \lambda.$$

Then we can generate Gomory's mixed-integer cut associated with this equation. Repeating this procedure for all every possible  $\lambda$ , we may apply all possible Gomory's mixed-integer cuts. We refer to the resulting set as the **mixed integer closure** of  $P$ .

**Theorem 17.3.** *The mixed integer closure of  $P$  coincides with the split closure of  $P$ .*

## References

- [CKS90] W.J. Cook, R. Kannan, and A. Schrijver. Chvátal closures for mixed integer programming problems. *Mathematical Programming*, 47:155–174, 1990. [17.1](#)