## 1 Outline

In this lecture, we study

- Minkowski-Weyl theorem for general polyhedra.
- recession cone and lineality space.
- implicit equalities and affine hull.

## 2 Minkowski-Weyl theorem for general polyhedra

A set  $P \subseteq \mathbb{R}^d$  is a **polyhedron** if it is defined by a **finite** number of linear inequalities, i.e.

$$P = \{ x \in \mathbb{R}^d : Ax \le b \}.$$

Hence, a polyhedron is a finite intersection of half-spaces. A polyhedron is rational if it is defined by

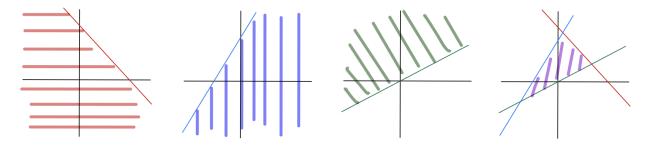


Figure 10.1: Polyhedron defined by three inequalities

a system of linear inequalities where all coefficients and right-hand sides are rational. A polyhedral cone is by definition a polyhedron.

Given a polyhedron  $P = \{x \in \mathbb{R}^d : Ax \leq b\}$ , we can associate a polyhedral cone given by

$$C_P = \left\{ (x, y) \in \mathbb{R}^d \times \mathbb{R} : Ax - by \le 0, y \ge 0 \right\}.$$

Then we have

$$P = \left\{ x \in \mathbb{R}^d : (x, 1) \in C_P \right\}.$$

Given two sets  $Q,C\subseteq \mathbb{R}^d,$  Minkowski sum of Q and C is defined as

$$Q + C = \left\{ x = u + v \in \mathbb{R}^d : u \in Q \ v \in C \right\}.$$

Based on these, we may prove the following theorem for polyhedra.

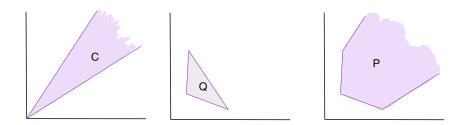


Figure 10.2: Polyhedron as the Minkowski sum of a polytope and a polyhedral cone

**Theorem 10.1** (Minkowski-Weyl theorem). A set  $P \subseteq \mathbb{R}^d$  is a polyhedron if and only if

$$P = \operatorname{conv}(v^1, \dots, v^p) + \operatorname{cone}(r^1, \dots, r^q)$$

for some vectors  $v^1, \ldots, v^p$  and  $r^1, \ldots, r^q$ .

*Proof.* ( $\Rightarrow$ ) Let  $P = \{x : \mathbb{R}^d : Ax \leq b\}$  be a polyhedron. Then consider the associated cone  $C_P = \{(x, y) \in \mathbb{R}^d \times \mathbb{R} : Ax - by \leq 0, y \geq 0\}$ . By Minkowski-Weyl theorem for cones, there exist vectors  $(u^1, w^1), \ldots, (u^k, w^k) \in \mathbb{R}^d \times \mathbb{R}$  such that

$$C_P = \operatorname{cone}\left(\binom{u^1}{w^1}, \dots, \binom{u^k}{w^k}\right).$$

By the definition of  $C_P$ , we have  $w^1, \ldots, w^k \ge 0$ . Then each  $w^i$  is strictly positive or equal to 0. For each vector  $(u^i, w^i)$  with  $w^i > 0$ , we divide it by  $w^i$ . Then  $C_P$  can be rewritten as

$$C_P = \operatorname{cone}\left(\binom{v^1}{1}, \dots, \binom{v^p}{1}, \binom{r^1}{0}, \dots, \binom{r^q}{0}\right)$$
$$= \left\{ (x, y) \in \mathbb{R}^d \times \mathbb{R} : \exists \lambda \ge 0, \mu \ge 0 \text{ s.t. } x = \sum_{i=1}^p \lambda_i v^i + \sum_{j=1}^q \mu_j r^j, \ y = \sum_{i=1}^p \lambda_i \right\}$$

Since  $P = \{x \in \mathbb{R}^d : (x, 1) \in C_P\}$ , it follows that

$$P = \operatorname{conv}(v^1, \dots, v^p) + \operatorname{cone}(r^1, \dots, r^q).$$

 $(\Leftarrow)$  Let  $C_P$  be defined as

$$C_P = \operatorname{cone}\left(\binom{v^1}{1}, \dots, \binom{v^p}{1}, \binom{r^1}{0}, \dots, \binom{r^q}{0}\right)$$

By Minkowski-Weyl theorem for cones, we know that  $C_P$  is a polyhedral cone, and therefore,  $C_P$  can be written as

$$C_P = \left\{ (x, y) \in \mathbb{R}^d \times \mathbb{R} : Ax - by \le 0 \right\}.$$

Note that  $P = \{x \in \mathbb{R}^d : (x, 1) \in C_P\}$ , in which case, we have

$$P = \{ x \in \mathbb{R}^d : Ax \le b \},\$$

as required.

A set  $P \subseteq \mathbb{R}^d$  is a **polytope** if it is a polyhedron and bounded, i.e.,  $P \subseteq [-M, M]^d$  for some sufficiently large M > 0.

**Corollary 10.2** (Minkowski-Weyl theorem for polytopes). A set  $P \subseteq \mathbb{R}^d$  is a polytope if and only if

$$P = \operatorname{conv}(v^1, \dots, v^p)$$

for some vectors  $v^1, \ldots, v^p$ .

*Proof.* ( $\Rightarrow$ ) By Theorem 10.1, as P is a polyhedron,  $P = \operatorname{conv}(v^1, \ldots, v^p) + \operatorname{cone}(r^1, \ldots, r^q)$  for some vectors  $v^1, \ldots, v^p$  and  $r^1, \ldots, r^q$ . Here, since P is bounded, we have q = 0 or  $r^1 = \cdots = r^q = 0$ , in which case  $P = \operatorname{conv}(v^1, \ldots, v^p)$ .

( $\Leftarrow$ ) By Theorem 10.1, P is a polyhedron. Moreover, since P is the convex hull of a finite number of vectors, P is bounded. Therefore, P is a polytope.

## **3** Recession cone and lineality space

Given a nonempty polyhedron P, a ray of P is a vector r such that

 $x + \lambda r \in P$  for all  $x \in P$  and  $\lambda \ge 0$ .

The **recession cone** of P is the set of all rays of P, i.e.,

$$\operatorname{rec}(P) = \left\{ r \in \mathbb{R}^d : x + \lambda r \in P \quad \text{for all } x \in P \text{ and } \lambda \ge 0 \right\}.$$

**Lemma 10.3.** rec(P) is a convex cone.

*Proof.* Let  $r^1, r^2 \in rec(P)$ . Then  $x + \lambda r^1, x + \lambda r^2 \in P$  for any  $x \in P$  and  $\lambda \ge 0$ . Since P is convex, for any  $0 \le \alpha \le 1$ , we have

$$x + \lambda \left( \alpha r^1 + (1 - \alpha)r^2 \right) = \alpha (x + \lambda r^1) + (1 - \alpha)(x + \lambda r^2) \in P.$$

This implies that  $\alpha r^1 + (1 - \alpha)r^2 \in rec(P)$ , and therefore, rec(P) is a convex cone.

A ray r of polyhedron P is called an **extreme ray** of P if it cannot be written as a convex combination of two distinct rays of P.

The **lineality space** of P is defined as

$$lin(P) = \left\{ r \in \mathbb{R}^d : x + \lambda r \in P \text{ for all } x \in P \text{ and } \lambda \in \mathbb{R} \right\}.$$

Hence, the lineality space is the set of all vectors r such that both r and -r are rays of P.

**Lemma 10.4.**  $lin(P) = rec(P) \cap - rec(P)$ .

*Proof.* Note that  $r \in \text{lin}(P)$  if and only if  $r \in \text{rec}(P)$  and  $-r \in \text{rec}(P)$ . Moreover,  $-r \in \text{rec}(P)$  if and only if  $r \in -\text{rec}(P)$ . Therefore,  $r \in \text{lin}(P)$  if and only if  $r \in \text{rec}(P)$  and  $r \in -\text{rec}(P)$ .  $\Box$ 

Moreover,

**Lemma 10.5.** lin(P) is a linear subspace.

*Proof.* Let  $r^1, r^2 \in \text{lin}(P)$ . Then  $x + 2\lambda r^1, x + 2\lambda r^2 \in P$  for any  $x \in P$  and  $\lambda \in \mathbb{R}$ . Since P is convex, we have

$$x + \lambda (r^1 + r^2) = \frac{1}{2}(x + 2\lambda r^1) + \frac{1}{2}(x + 2\lambda r^2) \in P.$$

This implies that  $r^1 + r^2 \in \text{lin}(P)$ . Moreover, for any  $\alpha \in \mathbb{R}$ , we know that  $x + \lambda \alpha r^1 \in P$ . This implies that  $\alpha r^1 \in \text{lin}(P)$ . Therefore, lin(P) is a linear subspace.

We say that polyhedron P is **pointed** if its lineality space is **trivial**, i.e.,  $lin(P) = \{0\}$ .

**Proposition 10.6.** Let  $P = \{x \in \mathbb{R}^d : Ax \leq b\}$  be a polyhedron such that

 $P = \operatorname{conv}(v^1, \dots, v^p) + \operatorname{cone}(r^1, \dots, r^q)$ 

for some vectors  $v^1, \ldots, v^p$  and  $r^1, \ldots, r^{q^1}$ . Then the recession cone of P is given by

$$\operatorname{rec}(P) = \left\{ x \in \mathbb{R}^d : Ax \le 0 \right\} = \operatorname{cone}(r^1, \dots, r^q).$$

Moreover, the lineality space is given by

$$\ln(P) = \left\{ x \in \mathbb{R}^d : Ax = 0 \right\} = \operatorname{cone}(r^1, \dots, r^q) \cap \operatorname{cone}(-r^1, \dots, -r^q).$$

*Proof.* We first show that  $\operatorname{rec}(P) = \{x \in \mathbb{R}^d : Ax \leq 0\}$ . Let r satify  $Ar \leq 0$ . Then for any  $x \in P$  and  $\lambda \geq 0$ , we have  $A(x + \lambda r) \leq Ax \leq b$ , implying that  $x + \lambda r \in P$ . Hence  $r \in \operatorname{rec}(P)$ . Let r be such that  $Ar \leq 0$ . Then  $(Ar)_i > 0$  for some component i. Then for any  $x \in P$ , there exists a sufficiently large  $\lambda > 0$  such that  $(Ax)_i + \lambda(Ar)_i > 0$ . In this case,  $Ax + \lambda Ar = A(x + \lambda r) \leq 0$ , and therefore,  $r \notin \operatorname{rec}(P)$ .

Next we show that  $\operatorname{rec}(P) = \operatorname{cone}(r^1, \ldots, r^q)$ . If  $r \in \operatorname{cone}(r^1, \ldots, r^q)$ , then  $x + \lambda r \in P$  for any  $x \in P$  and  $\lambda \geq 0$ . Therefore,  $r \in \operatorname{rec}(P)$ . Let  $r \in \operatorname{rec}(P)$ . Then for any  $x \in \operatorname{conv}(v^1, \ldots, v^p)$ , we have  $x + \lambda r \in P$  for  $\lambda \geq 0$ . Since  $\operatorname{conv}(v^1, \ldots, v^p)$  is bounded, we must have  $r \in \operatorname{cone}(r^1, \ldots, r^q)$ .

By the previous lemma, we have that  $lin(P) = rec(P) \cap - rec(P)$ . Note that

$$-\operatorname{rec}(P) = \left\{ x \in \mathbb{R}^d : Ax \ge 0 \right\} = \operatorname{cone}(-r^1, \dots, -r^q).$$

as required.

## 4 Implicit equalities and affine hull

Consider a polyhedron  $P = \{x \in \mathbb{R}^d : Ax \leq b\}$  where  $Ax \leq b$  consists of linear inequalities  $a_i^{\top}x \leq b_i$  for  $i \in [m]$ . We say that  $a_i^{\top}x \leq b_i$  is an **implicit equality** if

$$P \subseteq \left\{ x \in \mathbb{R}^d : a_i^\top x = b_i \right\}.$$

In words,  $a_i^{\top} x \leq b_i$  is an implicit equality if every point in P satisfies it with equality. If there is a point x in P such that  $a_i^{\top} < b_i$ , i.e., x satisfies the inequality strictly, then  $a_i^{\top} x \leq b_i$  is not an implicit equality.

Let  $A^{=}x \leq b^{=}$  be the subsystem of  $Ax \leq b$  that collects all implicit equalities, and let  $A^{<}x \leq b^{<}$  collect the other inequalities in  $Ax \leq b$ .

<sup>&</sup>lt;sup>1</sup>These vectors exist due to Minkowski-Weyl theorem

**Lemma 10.7.** There is a point  $\bar{x} \in P$  that satisfies  $A^{<}\bar{x} < b^{<}$ , i.e.,  $\bar{x}$  satisfies all inequalities  $A^{<}x \leq b^{<}$  strictly.

*Proof.* Suppose that  $A^{\leq}x \leq b^{\leq}$  is given by  $a_i^{\top}x \leq b_i$  for  $i \in I$  where I is some subset of [m]. For each  $i \in I$ , as  $a_i^{\top}x \leq b_i$  is not an implicit equality, there exists some  $x^i \in P$  such that  $a_i^{\top}x^i < b_i$ . Morever, as  $x^i$  is point in P, it satisfies  $Ax^i \leq b$ . Then we take a convex combination of  $x^i$  for  $i \in I$  given by

$$\bar{x} = \frac{1}{|I|} \sum_{i \in I} x^i.$$

Then for any  $i \in I$ , we have  $a_i^\top \bar{x} < b_i$ .

**Theorem 10.8.** The affine hull of polyhedron P is given by

aff(P) = 
$$\left\{ x \in \mathbb{R}^d : A^= x = b^= \right\} = \left\{ x \in \mathbb{R}^d : A^= x \le b^= \right\}.$$

In particular,  $\dim(P) = d - \operatorname{rank}(A^{=})$ .

*Proof.* It is straightforward that

aff
$$(P) \subseteq \left\{ x \in \mathbb{R}^d : A^= x = b^= \right\} \subseteq \left\{ x \in \mathbb{R}^d : A^= x \le b^= \right\}$$

Then it suffices to argue that

$$\left\{x \in \mathbb{R}^d : A^{=}x \le b^{=}\right\} \subseteq \operatorname{aff}(P)$$

Let  $\hat{x}$  satisfy  $A^{=}\hat{x} \leq b^{=}$ . By Lemma 10.7, there exists  $\bar{x} \in P$  such that  $A^{<}\bar{x} < b^{<}$ . Then for some sufficiently small  $\epsilon > 0$ , we have

$$A^{<}(\bar{x} + \epsilon(\hat{x} - \bar{x})) = A^{<}\bar{x} + \epsilon A^{<}(\hat{x} - \bar{x}) \le b^{<}.$$

Moreover,

$$A^{=}(\bar{x} + \epsilon(\hat{x} - \bar{x})) = (1 - \epsilon)A^{=}\bar{x} + \epsilon A^{=}\hat{x} \le b^{=}$$

Let  $\tilde{x} = \bar{x} + \epsilon(\hat{x} - \bar{x})$ . Then it follows that  $\tilde{x} \in P$ . Note that as  $\bar{x}, \tilde{x} \in P$ , the line going through  $\bar{x}$  and  $\tilde{x}$  is contained in the affine hull aff(P). Moreover,

$$\hat{x} = \frac{1}{\epsilon}\tilde{x} - \frac{1-\epsilon}{\epsilon}\hat{x}.$$

Here, the coefficients sum up to 1, and therefore,  $\hat{x}$  is on the line going through  $\bar{x}$  and  $\tilde{x}$ . Therefore,  $\hat{x} \in \operatorname{aff}(P)$ .

Polyhedron  $P = \{x \in \mathbb{R}^d : Ax \leq b\}$  is **full-dimensional** if dim(P) = d, in which case, the system  $Ax \leq b$  does not involve an implicit equality.

**Example 10.9.** Let us consider the **assignment polytope**, given by

$$P = \left\{ x \in \mathbb{R}^{n^2} : \sum_{i=1}^n x_{ij} = 1, \quad i = 1, \dots, n \\ \sum_{i=1}^n x_{ij} = 1, \quad j = 1, \dots, n \\ x_{ij} \ge 0, \quad i, j = 1, \dots, n \end{array} \right\}.$$

We may prove that  $\dim(P) = n^2 - 2n + 1$ .

Note that the system defining P has 2n equality constraints. Let Ax = 1 denote the system that collects the 2n equality constraints. We will show that rank(A) = 2n - 1. Note that

$$\sum_{i=1} \sum_{j=1} x_{ij} - \sum_{j=1} \sum_{i=1} x_{ij} = 0.$$
aggregating the first set of equalities aggregating the second set

This implies that A is not of full row rank, and therefore,  $\operatorname{rank}(A) \leq 2n - 1$ . Next, consider the column submatrix of A associated with variables  $x_{1i}$  for  $i = 1, \ldots, n$  and  $x_{ii}$  for  $i = 2, \ldots, n$ .

The submatrix has 2n - 1 columns which are linearly independent. Therefore,  $\operatorname{rank}(A) \ge 2n - 1$ . This implies that  $\operatorname{rank}(A) = 2n - 1$ . Then it follows from Theorem 10.8 that  $\dim(P) = n^2 - \operatorname{rank}(A) = n^2 - 2n + 1$ , as required.