1 Outline

In this lecture, we study

- complexity of nonconvex optimization,
- sparse regression,
- low-rank matrix completion.

2 Introduction to Nonconvex Optimization

Figure 9.1 shows a nonconvex function with two variables. As we can see, the function has one

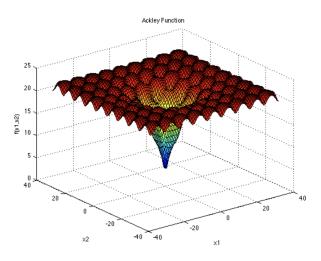


Figure 9.1: Ackley Function with Two Variables

global minimum but it has multiple local minima. Recall that for a convex function f, if $\nabla f(x) = 0$, then the corresponding x is a minimizer of the function f. Gradient-based methods for convex optimization basically seek for a solution x with $\nabla f(x) = 0$. However, for a nonconvex function f, a point x with $\nabla f(x) = 0$ does not guarantee global optimality as it can be a locally optimal solution.

In fact, finding the global minimum of a nonconvex function is a difficult task in general. To be more precise, the following theorem shows that there exists a smooth nonconvex function, to find the global minimum of which we need an exponential number of function evaluations.

Theorem 9.1. For any $\beta > 0$, there exists a β -smooth function $f : [0,1]^d \to \mathbb{R}$ on $[0,1]^d$ such that any algorithm requires at least $(\beta/\epsilon)^{\Omega(d)}$ function queries to find an ϵ -optimal solution x with

$$f(x) \le \min_{x \in \mathbb{R}^d} f(x) + \epsilon.$$

Proof. We partition [0, 1] into k invervals of equal length, which gives rise to k^d boxes partitioning $[0, 1]^d$. We can construct a function f such that a box contains a point x^* with $f(x^*) = -\epsilon$ but f has value 0 in the other boxes. This means that checking a box not containing x^* does not provide any information about the location of x^* . This means that to find the box containing x^* , we need at least $\Omega(k^d)$ function evalutions. To make function f smooth with parameter β , we can make f behaves like

$$f(x) \simeq f(x^*) + \frac{\beta}{2} ||x - x^*||_2^2$$

At the same time, to impose the condition that f(x) = 0 if x is not contained in the box with x^* , we can set $k = O(\sqrt{\beta/\epsilon})$. Therefore, we need at least $O\left((\sqrt{\beta/\epsilon})^d\right)$ function evaluations. \Box

There exist several important applications of nonconvex optimization. In the remainder of this section, we provide an overview of some well-known nonconvex optimization problems.

2.1 Sparse Regression

Let us consider the following optimization problem.

$$\min_{x \in \mathbb{R}^d} \quad \frac{1}{2} \|Ax - b\|_2^2 + \lambda \|x\|_0 \tag{9.1}$$

where

$$||x||_0 = |\{i \in [d] : x_i \neq 0\}|$$

counts the number of nonzero coordinates of x. Here, $||x||_0$ is called the ℓ_0 -norm. Although the name of $||x||_0$ contains the term "norm", $||x||_0$ is a nonconvex function and thus it is not a norm. The optimization problem is referred to as **sparse regression** because the $\lambda ||x||_0$ term encourages to use less variables of x.

Due to nonconvexity of $||x||_0$, it is often difficult to solve (9.1) efficiently. Motivated by this, we can approximate and replace the ℓ_0 -norm by the ℓ_1 -norm, which gives rise to LASSO:

$$\min_{x \in \mathbb{R}^d} \quad \frac{1}{2} \|Ax - b\|_2^2 + \lambda \|x\|_1.$$
(9.2)

We saw that the proximal gradient method solves (9.2) and it runs with

$$x_{t+1} = \operatorname{prox}_{\eta\lambda\|\cdot\|_1}(x_t - \eta\nabla f(x_t))$$

where

$$f(x) = \frac{1}{2} \|Ax - b\|_2^2 \quad \text{and} \quad \left(\operatorname{prox}_{\eta \lambda \| \cdot \|_1}(x) \right)_{i \in [d]} = \begin{cases} x - \eta \lambda, & \text{if } x \ge \eta \lambda, \\ 0, & \text{if } -\eta \lambda \le x < \eta \lambda \\ x + \eta \lambda, & \text{if } x < -\eta \lambda. \end{cases}$$

Recall that the prox operator $\operatorname{prox}_{\eta\lambda\|\cdot\|_1}(\cdot)$ is called the shrinkage operator or the **soft-thresholding** operator. In fact, we can apply proximal gradient descent to solve (9.1). Note that

$$\operatorname{prox}_{\eta\lambda\|\cdot\|_0}(x) = \operatorname{argmin}_{u\in\mathbb{R}^d} \left\{ \|u\|_0 + \frac{1}{2\eta\lambda} \|x-u\|_2^2 \right\}.$$

By definition,

$$\left(\operatorname{prox}_{\eta\lambda\|\cdot\|_{0}}(x)\right)_{i\in[d]} = \begin{cases} x_{i}, & \text{if } x_{i}^{2} \geq 2\eta\lambda, \\ 0, & \text{otherwise.} \end{cases}$$

Here, the operator $\operatorname{prox}_{\eta\lambda\|\cdot\|_0}(\cdot)$ is referred to as the **hard-thresholding** operator. With the hard-thresholding operator, one may run proximal gradient descent with update rule

$$x_{t+1} = \operatorname{prox}_{\eta \lambda \parallel \cdot \parallel_0} (x_t - \eta \nabla f(x_t)).$$

Unlike LASSO, however, proximal gradient descent applied to sparse regression does not necessarily converge to an optimal solution. This is again due to nonconvexity of the ℓ_0 -norm.

Recent works [HM20, HMS22] consider mixed-integer programming formulations of (9.1) while earlier works including LASSO focused on convex approximations of (9.1).

2.2 Low-Rank Matrix Completion

Let us consider

$$\min_{X \in \mathbb{R}^{n \times p}} \quad \|D - X\|_F \quad \text{subject to} \quad \operatorname{rank}(X) = k$$

where

- D is an $n \times p$ matrix,
- $||A||_F$ denotes the Frobenius norm, i.e., $||A||_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^p a_{ij}^2}$.

By the definition of the Frobenius norm, the problem is equivalent to

$$\min_{X \in \mathbb{R}^{n \times p}} \quad \frac{1}{2} \sum_{i=1}^{n} \sum_{j=1}^{p} (d_{ij} - x_{ij})^2 \quad \text{subject to} \quad \text{rank}(X) = k.$$

A common application is movie recommendation where matrix D collects user preference data of n users rating p movies. In movie recommendation, the user rating matrix D is typically very sparse, as it is rare that an user rates all movies. Hence, the goal is to find a matrix X that completes the missing entries of D. The general hypothesis is that the true rating matrix $X \in \mathbb{R}^{n \times p}$ is generated by the product of an user-feature matrix $U \in \mathbb{R}^{n \times k}$ and a movie-feature matrix $V \in \mathbb{R}^{p \times k}$ over k features as follows.

$$X \qquad \left] = \left[\begin{array}{c} U \\ U \end{array} \right] \left[\begin{array}{c} V^{\top} \\ \end{array} \right].$$

Under such a hypothesis, matrix X has rank k.

As expected, the constraint $\operatorname{rank}(X) = k$ defines a nonconvex set, so the matrix completion problem is nonconvex. A commonly used solution approach is based on relaxing the constraint $\operatorname{rank}(X) = k$ by

$$||X||_* = \operatorname{Trace}(\sqrt{X^{\top}X}) = \sum_{i=1}^{\min\{n,p\}} \sigma_i(X) \le k$$

where $||X||_*$ denotes the **nuclear norm** of X and $\sigma_1(X), \ldots, \sigma_{\min\{n,p\}}(X)$ are the singular values of X.

2.3 Max-Cut

Given a graph G = (V, E), the **max-cut** problem seeks to find a partition the vertex set V so that the number of edges crossing the partition is maximized. Here, a partition (V_1, V_2) of V consists of two sets V_1, V_2 satisfying $V_1 \cup V_2 = V$ and $V_1 \cap V_2 = \emptyset$, and the set of edges crossing the partition is basically $\{uv \in E : u \in V_1, v \in V_2\}$. For example, in Figure 9.2, there is a graph of 5 vertices partitioned into red and black vertices, and the edges highlighted are the ones crossing the partition.

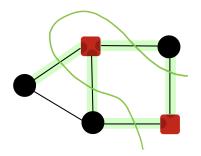


Figure 9.2: Edges crossing a partition

The problem can be formulated by the following (discrete) optimization problem:

maximize
$$\sum_{ij\in E} \frac{1-x_i x_j}{2}$$

subject to $x_i \in \{-1,1\}$ for $i \in V$.

As long as $x_i \in \mathbb{R}, x_i \in \{-1, 1\}$ is equivalent to $x_i^2 = 1$. Hence, the formulation is equivalent to

maximize
$$\sum_{ij \in E} \frac{1 - x_i x_j}{2}$$

subject to $x_i^2 = 1$ for $i \in V$.

Note that the contstraint $x_i \in \{-1, 1\}$ as well as $x_i^2 = 1$ are nonconvex constraints. The vector relaxation of the formulation is obtained by replacing x_i by vector $v_i \in \mathbb{R}^k$ as follows.

maximize
$$\sum_{ij \in E} \frac{1 - v_i^\top v_j}{2}$$

subject to $||v_i||_2 = 1$ for $i \in V$.

Again, the constraint $||v_i||_2 = 1$ is still nonconvex.

Another relaxation is given as follows. Let d = |V|. Then we consider a $d \times d$ matrix X whose entry at *i*th row and *j*th column, X_{ij} , is $x_i x_j$. Then we have that $X = xx^{\top}$, which is the outer product of vector x and itself. In fact, X is of the form $X = xx^{\top}$ if and only if X is positive semidefinite and

the rank of X is precise 1. What this implies is that, the max-cut formulation can be rewritten as

maximize
$$\sum_{ij\in E} \frac{1-X_{ij}}{2}$$
subject to $X_{ii} = 1$ for $i \in V$,
 $X \succeq 0$,
rank $(X) = 1$.

Here, the constraint $\operatorname{rank}(X) = 1$ is nonconvex. A common approach is to take out the nonconvex constraint and consider

maximize
$$\sum_{ij\in E} \frac{1-X_{ij}}{2}$$

subject to $X_{ii} = 1$ for $i \in V$,
 $X \succeq 0$.

This is often called the **semidefinite programming (SDP) relaxation** of max-cut. Here, the SDP relaxation is a convex optimization problem.

References

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- [HMS22] Hussein Hazimeh, Rahul Mazumder, and Ali Saab. Sparse regression at scale: branchand-bound rooted in first-order optimization. *Mathematical Programming*, page 347–388, 2022. 2.1